

ASSIGNMENT 6 SOLUTIONS

MATH 303, FALL 2011

If you find any errors please let me know

MANIPULATION

(M1) By rule C we have that

$$((c = d) \wedge (d = c')) \rightarrow c = c'$$

is valid. Then by rule F applied with $A(x)$ being $((c = x) \wedge (x = c'))$ and B being $c = c'$ we get that

$$\exists x((c = x) \wedge (x = c')) \rightarrow c = c'$$

is valid, which is what we were supposed to show.

(M2) First notice that by Cohen's definition of "derive" the question is asking if

$$(\forall xA(x) \wedge (A(c) \rightarrow B)) \rightarrow B$$

is valid.

By rule E we know that $\forall xA(x) \rightarrow A(c)$ is valid. Consider the following formula

$$(\forall xA(x) \rightarrow A(c)) \rightarrow ((\forall xA(x) \wedge (A(c) \rightarrow B)) \rightarrow B)$$

Letting C be $\forall xA(x)$ this is the formula

$$(C \rightarrow A(c)) \rightarrow ((C \wedge (A(c) \rightarrow B)) \rightarrow B)$$

This is a propositional function in the letters C , $A(c)$, and B , so we can apply rule A. To do this build a truth table (unfortunately a bit of a large one)

$A(c)$	B	C	$C \rightarrow A(c)$	$A(c) \rightarrow B$	$C \wedge (A(c) \rightarrow B)$	$(C \wedge (A(c) \rightarrow B)) \rightarrow B$	whole thing
F	F	F	T	T	F	T	T
F	F	T	F	T	T	F	T
F	T	F	T	T	F	T	T
F	T	T	F	T	T	T	T
T	F	F	T	F	F	T	T
T	F	T	T	F	F	T	T
T	T	F	T	T	F	T	T
T	T	T	T	T	T	T	T

We see that the propositional function in question is identically true, so by rule A it is a valid statement. Finally since $CA(c)$ is valid and $(C \rightarrow A(c)) \rightarrow ((C \wedge (A(c) \rightarrow B)) \rightarrow B)$ by rule B we conclude that

$$(C \wedge (A(c) \rightarrow B)) \rightarrow B$$

which is what we wanted to show.

- (M3) $\{a, b, c, d\}$ is both maximal and largest in X . $\{a\}$, $\{b\}$, $\{c\}$, and $\{d\}$ are all minimal, but there is no smallest element (if there were it would also have to be the unique minimal element, but there are 4 minimal elements).
- (M4) $((\omega^+)^+)^+ = \{0, 1, 2, \dots, \omega, \omega^+, (\omega^+)^+\}$ Take any two elements m and n of $((\omega^+)^+)^+$, I will describe precisely when $m \leq n$ in $((\omega^+)^+)^+$.
 If m and n are equal then $m \leq n$ in $((\omega^+)^+)^+$.
 If m and n are both natural numbers then $m \leq n$ in $((\omega^+)^+)^+$ if $m \leq n$ in ω .
 If m is a natural number and n is not then $m \leq n$. (If n is a natural number and m is not then $m \not\leq n$.)
 If neither m and n are natural numbers but they are distinct, then set $\omega \leq \omega^+$, $\omega^+ \leq (\omega^+)^+$ (and hence $\omega \leq (\omega^+)^+$) to determine when $m \leq n$.
- (M5) $\omega 3 = \{0, 1, 2, 3, \dots, \omega, \omega + 1, \omega + 2, \dots, \omega 2, \omega 2 + 1, \omega 2 + 2, \dots\}$. Define $f : \omega 3 \rightarrow \omega$ as follows.

$$f(x) = \begin{cases} 3x & \text{if } x \in \omega \\ 3k + 1 & \text{if } x = \omega + k \\ 3k + 2 & \text{if } x = \omega 2 + k \end{cases}$$

Now we just need to show that f is one-to-one and onto.

One-to-one: Suppose $f(a) = f(b)$. If $f(a)$ is divisible by 3 we have $3a = 3b$ so $a = b$. If $f(a)$ is congruent to 1 modulo 3 then we have $3k + 1 = 3\ell + 1$ where $a = \omega + k$ and $b = \omega + \ell$, so $k = \ell$ and so $a = b$. Similarly if $f(a)$ is congruent to 2 modulo 3 then we have $3k + 1 = 3\ell + 1$ where $a = \omega 2 + k$ and $b = \omega 3 + \ell$ so again $k = \ell$ and $a = b$.

Onto: Take any $n \in \omega$. If n is divisible by 3 then $n = 3x$ and $f(x) = n$. If $n = 3k + 1$ then $f(\omega + k) = n$. If $n = 3k + 2$ then $f(\omega 2 + k) = n$. This exhausts the possibilities so f is onto.

PURE MATH

- (P1) (a) a a least element of X means that for all $x \in X$, $a \leq x$. Suppose $b \in X$ with $b \leq a$, then since a is least we also have $a \leq b$ and so $a = b$. Thus a is minimal.
 (b) Suppose a and b are both least elements of X . Then $a \leq b$ and $b \leq a$, so $a = b$. Thus if X has a least element then it has a unique least element.
- (P2) (a) Let R be addition, that is $(a, b, c) \in \bar{R}$ if and only if $a + b = c$. Then the sentence of S says $\forall x \forall y \forall z (\forall w (x + y = w \leftrightarrow x + z = w) \rightarrow y = z)$ equivalently $\forall x \forall y \forall z (x + y = x + z \rightarrow y = z)$. This is true in \mathbb{Z} (just subtract x from both sides). Thus with this interpretation \mathbb{Z} is a model for S .
 (b) *most possible choices for \bar{R} will make S false, here is one possibility:* Let R be the statement that the first and last arguments agree, that is $(a, b, c) \in \bar{R}$ if and only if $a = c$. Then the sentence of S says $\forall x \forall y \forall z (\forall w (x = w \leftrightarrow x = w) \rightarrow (y = z))$. $x = w \leftrightarrow x = w$ is always true, so this is just the sentence $\forall x \forall y \forall z \forall w (y = z)$, equivalently $\forall y \forall z (y = z)$ which is false in \mathbb{Z} since \mathbb{Z} has at least two elements. Thus with this interpretation \mathbb{Z} is not a model for S .

