

① Review of well orders

Recall that if X with \leq is a partially ordered set
and

④ for all $a, b \in X$ $a \leq b$ or $b \leq a$

⑤ for my ^{nonempty} $Y \subseteq X$, Y has a least element.
ie. there is an $a \in Y$ such that $a \leq x$ for all $x \in Y$.

Then we say \leq is a **well order** and we
say X is **well ordered** by \leq

Note ④ isn't needed because ⑤ implies ④

take $a, b \in X$ let $Y = \{a, b\}$

Then by ⑤ Y has a least element

If the least element is a then $a \leq b$

If the least element is b then $b \leq a$

so you get either $a \leq b$ or $b \leq a$.

Note if X with \leq is a partial order and satisfies (4)
then \leq is a total order on X

We saw that ω is well ordered by the usual \leq

eg Consider $\omega \times \omega$

What ordering shall we use?

(a) say $(a, b) \leq (c, d)$ if $a \leq c$ and $b \leq d$

eg $(1, 256) \leq (3, 8000000)$

$(1, 256) \neq (3, 4)$

Is $\omega \times \omega$ well ordered?

No. eg $\{(1, 256), (3, 4)\}$
has no least element
since $(1, 256) \neq (3, 4)$
and $(3, 4) \neq (1, 256)$

(b) say $(a, b) \leq (c, d)$ if
 $a \leq c$
or $(a = c \text{ and } b \leq d)$

This is called **lexicographic order** because it is how we put things in alphabetical order.

Is $\omega \times \omega$ well ordered?

Yes! let $\gamma \subseteq \omega \times \omega$

let $Z \subseteq \gamma$ be the set of
 $(a, b) \in \gamma$ where a is
smallest among all
first coordinates of elements
of γ

$\therefore Z = \{(a, b), (a, b_1), (a, b_2), \dots\}$

since ω is well ordered the
first coordinates of elements of γ
are well ordered so such
an a exists.

Furthermore \mathbb{Z} has a least element because all elements of \mathbb{Z} have the same first coordinate and the second coordinates are well ordered.

Let (a, b) be the least element of \mathbb{Z} .

All elements of \mathbb{Y} are at least as big so (a, b) is the least element of \mathbb{Y} .

② Transfinite induction

Well ordered sets have the nice property that we have a notion of induction. First we need the following definition

Definition

Let X with \leq be a partially ordered set.

Take $a \in X$. Then set

$$s(a) = \{x \in X : x \leq a\}$$

is called the **initial segment** of a

eg $X = \omega$ $a = 12$, $s(a) = \{0, 1, 2, \dots, 11\} \stackrel{\text{as it turns out}}{=} 12$

eg $E = \{a, b, c\}$, $X = P(E)$ ordered by \subseteq ,
 $s(\{a, b\}) = \{\emptyset, \{a\}, \{b\}\}$

let X be a well ordered set and let $S \subseteq X$

if for any $x \in X$ it is the case that

$s(x) \subseteq S$ implies $x \in S$

Then $S = X$

This is called the principle of transfinite induction

First lets see that this is true and then see what we can do with it

To check the fact suppose S has the property
but $S \neq X$.

then $X - S$ is nonempty so since X is well ordered $X - S$ has a least element,
call it a .

Consider $s(a)$ every $x \in s(a)$ has $x < a$

so by minimality of a , $x \in S$ for all $x \in s(a)$

Thus $s(a) \subseteq S$, so $a \in S$, contradicts $a \in X - S$

so $X - S$ can't be nonempty

so $X = S$

How does this relate to the principle of mathematical induction which we have already seen?

let $X = \omega$

First

Take $0 \in \omega$

$$s(0) = \emptyset$$

so transfinite induction says

$$s(0) \subseteq S \text{ implies } 0 \in S$$

but $\emptyset \subseteq S$ for any S

so if the transfinite induction property holds for $S \subseteq \omega$ then

Next note for $X = \omega$ transfinite induction ~~automatically~~ $0 \in S$

becomes strong induction, that is to conclude
 $x \in S$ you need the entire initial segment
inside S . I.e. your property needs to hold
for all $a < x$.

On the other hand the principle of mathematical induction is weak induction, that is, to conclude

$x \in S$ you only need $(x-1) \in S$

For ω strong and weak induction are equally powerful, they can prove the same statements.

But for other well ordered sets transfinite induction is necessary

e.g. let $X = \omega^+ = \omega \cup \{\omega\} = \{0, 1, 2, \dots, \omega\}$

use \leq given by, for $m, n \in \omega$

let $m \leq n$ in ω^+ if $m \leq n$ in ω
and vice versa

for $m \in \omega$, let $m \leq \omega$.

This is a well order.

Suppose we try to use the old principle of mathematical induction on X . What goes wrong?

For the break

Can you find an $S \subsetneq \omega^+$

with $0 \in S$ and for all $n \in S$, $n^+ \in S$?

answer ω The usual principle of mathematical induction is not strong enough to distinguish ω from ω^+

Transfinite induction fixes this problem

$$s(x) \subseteq S \Rightarrow x \in S$$

for ω : what is $s(\omega)$ in ω^+
 $\{0, 1, 2, 3, \dots\}$

so $s(\omega) = \omega$ so certainly $s(\omega) \subseteq \omega$
but $\omega \notin \omega$
and so transfinite induction
distinguishes ω from ω^+

One more note: transfinite induction also allows us to take base cases other than 0 without reindexing

How? just use the well ordered set $\{1, 2, \dots\}$
(or $\{5, 6, \dots\}$ or wherever you want).

③ Ordinals

We had

$$0, 0^+ = 1, 1^+ = 2, 2^+ = 3, 4, 5, \dots$$

all together we have

$$\{0, 1, 2, 3, \dots\} = \omega$$

Now consider

$$\omega, \omega^+, (\omega^+)^+, ((\omega^+)^+)^+, (((\omega^+)^+)^+)^+, \dots$$

what's next? Idea what is the next number after this, in the same way ω was after all natural numbers.

It should be $\{0, 1, 2, \dots, \omega, \omega^+, (\omega^+)^+, \dots\}$ but is this a set?

Suppose f is a function with domain $n \in \omega$

Say f is an ω -successor function

$$\text{if } f(0) = \omega \\ \text{and } f(m^+) = (f(m))^+$$

eg $n=3 = \{0, 1, 2\}$. $f(0) = \omega$
 $f(1) = \omega^+$ $f(1) = f(0^+) = (f(0))^+$
 $f(2) = (\omega^+)^+$ $= \omega^+$

then f is an ω -successor function.

In fact for each n there is a unique
 ω -successor function. Intuition we never had a chance
to make a choice in defining f , so it must have been unique.

Suppose f and g were both ω -successor functions
with domain n .

- $f(0) = \omega = g(0)$ so they agree on 0
- let $i \in n$ be the smallest number for which
 $f(i) \neq g(i)$.
using the well ordering of ω
also need it to start at 0
- $i \neq 0$ so $i = j^+$ for some $j \in \omega$ ($j = i - 1$)
but i was the smallest number where they
disagreed so $f(j) = g(j)$
so $f(i) = f(j^+) = f(j)^+ = g(j)^+ = g(j^+) = g(i)$
contradicting $f(i) \neq g(i)$.

Thus f is unique.

What we want is to join all those things together
ie show $\{0, 1, \dots, \omega, \omega^+, (\omega^+)^+, \dots\}$ is a set

let $S(n, x)$ be the property

" $n \in \omega$ and x is in the range of an
 ω -successor function with domain n "

in logic if can be done

but lets not bother

7 vs 2

The set we are looking for is

$$\{x : \exists n (\text{new} \wedge S(n, x))\}$$

We only know this is a set by the axiom of replacement

Intuitively S is acting like a function

it takes n to the range of the unique w -successor function defined on n

Call this function $F(n) = \{x : S(n, x)\}$

We want to know

- either (a) F is a function in the set theoretic sense (i.e. can write the set of ordered pairs defining F)
or (b) The image of any $X \subseteq w$ under F is a set.

These are equivalent. If F is a set theoretic function then its range is a set and can pull out ranges of subsets by the axiom of subset selection

If the images are sets, then the range of F , γ is a set and so have $w \times \gamma$ and can pull out F using the axiom of subset selection

- (b) would come from Cohen's version of the axiom of replacement
- (a) is Halmos' version which he calls the axiom of substitution

Axiom of Substitution If $S(a, b)$ is a sentence such that for each $a \in A$ the set $\{b : S(a, b)\}$ can be formed, then there exists a function F with domain A such that $F(a) = \{b : S(a, b)\}$ for each $a \in A$

Our use of the axiom of substitution will be, as above to extend our ability to count beyond w, w^t, \dots

So we have

$$0, 1, 2, 3, \dots$$

$$\omega, \omega^+, (\omega^+)^+, ((\omega^+)^+)^+, \dots \quad (*)$$

and by the above we can define a set theoretic function
 F with domain ω such that

$$F(0) = \omega, F(n^+) = (F(n))^+$$

let X be the range of F then $X = \{\omega, \omega^+, (\omega^+)^+, \dots\}$

Then the next number after the ones in $(*)$ is

$$X \cup \omega = \{0, 1, 2, \dots, \omega, \omega^+, (\omega^+)^+, \dots\}$$

So this really is a set

Note we went to all that work just to show
 $X \cup \omega$ is a set

What are these new bigger counting "numbers"
ordinals

Definition An ordinal is a well ordered set S
such that for all $x \in S$ $s(x) = x$

e.g. let's check 3 is an ordinal

$3 = \overbrace{\{0, 1, 2\}}$ ordered in the usual way $0 \leq 1 \leq 2$

$$s(0) = \emptyset = 0$$

$$s(1) = \{0\} = 1$$

$$s(2) = \{0, 1\} = 2$$

Likewise every natural number is an ordinal.

eg check ω is an ordinal.

Two useful facts

① If X is an ordinal then X^+ is an ordinal

proof Use the order on X^+ given by

This is a well order as

Finally we can check the ordinal property.

② Let X be a set. There is at most one well order which makes X into an ordinal

proof

Suppose there is a well order which makes X into an ordinal. Take any other well order of X

④

Next time

More on ordinals

Please read Halmos section 20