

Math 303, Fall 2011, Lecture 7

① Functions

Suppose X and Y are sets. We would like to encode functions

$f: X \rightarrow Y$ using sets

How can we do this?

name of function \nearrow f \nwarrow *domain* X \rightarrow Y \nwarrow *codomain*

Use ordered pairs

Just let f be the set of (x, y)
where $f(x) = y$

To make this a precise definition we need

$$f = \{ (x, y) \in ?? \mid ?? \}$$

and so $f(x) = y$ must be writable
in logic

Which subsets S of $X \times Y$ are functions?

for each $x \in X$ we need
a unique $y \in Y$ so $(x, y) \in S$

(Need at least one as otherwise
 $f(x)$ is not defined, and
no more than one since $f(x)$
can only have one value)

Define $Y^X = \{ f \in \mathcal{P}(X \times Y) \mid f \text{ is a function} \}$

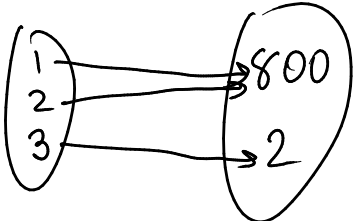
eg let $X = Y = \omega (= \{0, 1, 2, \dots\})$

let $f(x) = x^2$

What is f as a set?

Answer $\{(0, 0), (1, 1), (2, 4), (3, 9), (4, 16), \dots\}$

eg let $X = \{1, 2, 3\}$ $Y = \{800, 2\}$

let f be  What is f as a set?

Answer $\{(1, 800), (2, 800), (3, 2)\}$

eg let $X = \emptyset$, let $Y = \omega$

What can f be?

Answer $f = \emptyset$ the function does nothing since its domain is empty

eg What is \mathcal{Y}^\emptyset for any set \mathcal{Y}
Answer \emptyset

eg let $X = \{1, 2\}$ let $\mathcal{Y} = \{3, 4\}$
what is \mathcal{Y}^X ?

Answer

$1 \rightarrow 3$	$1 \rightarrow 4$	$1 \rightarrow 3$	$1 \rightarrow 4$
$2 \rightarrow 3$	$2 \rightarrow 4$	$2 \rightarrow 4$	$2 \rightarrow 3$

$\mathcal{Y}^X = \left\{ \left\{ (1,3), (2,3) \right\}, \left\{ (1,4), (2,4) \right\}, \left\{ (1,3), (2,4) \right\}, \left\{ (1,4), (2,3) \right\} \right\}$

lets remember some function words

domain $\text{dom } f = \{x : \text{for some } y \ (x, y) \in f\}$

range $\text{ran } f = \{y : \text{for some } x \ (x, y) \in f\}$

If $\text{ran } f = Y$ then f is **onto** Y

If $X \subseteq Y$ then the function $f: X \rightarrow Y$
defined by $f(x) = x$ for all $x \in X$
is called the **inclusion map**

The inclusion map from X into X is called
the **identity map**

The function $f: X \times Y \rightarrow Y$ given by $f(x, y) = y$
is the **projection map** onto the second coordinate

we also have

$$f: X \times Y \rightarrow X \quad \text{given by } f(x, y) = x$$

the **projection map** onto the first coordinate

If f maps distinct elements onto distinct elements

then f is **one-to-one**

written in function language f one-to-one means

$$\text{if } f(x_1) = f(x_2) \quad \text{then } x_1 = x_2$$

written in set language f one-to-one means

$$\text{if } (x_1, y) \in f \quad \text{and } (x_2, y) \in f$$

$$\text{then } x_1 = x_2$$

eg let $X = \{1, 2\}$ $Y = \{1, 3\}$

What is projection from $X \times Y$ to Y ?

Answer $f(x, y) = y$ for all $(x, y) \in X \times Y$

As a set

$\{ \{(1, 1), 1\}, \{(1, 3), 3\}, \{(2, 1), 1\}, \{(2, 3), 3\} \}$

eg is the projection in the previous example onto Y ?

Answer, yes

is it one-to-one?

Answer, no

② Finite and infinite (This is back to Halmos p52)

Definition Say two sets A and B are **equivalent** or **the same size** if there is a **one-to-one and onto** function from A to B .

could mean other things in other books so watch out.

eg is $\{2, 3, 8, 12\}$ equivalent to 4?

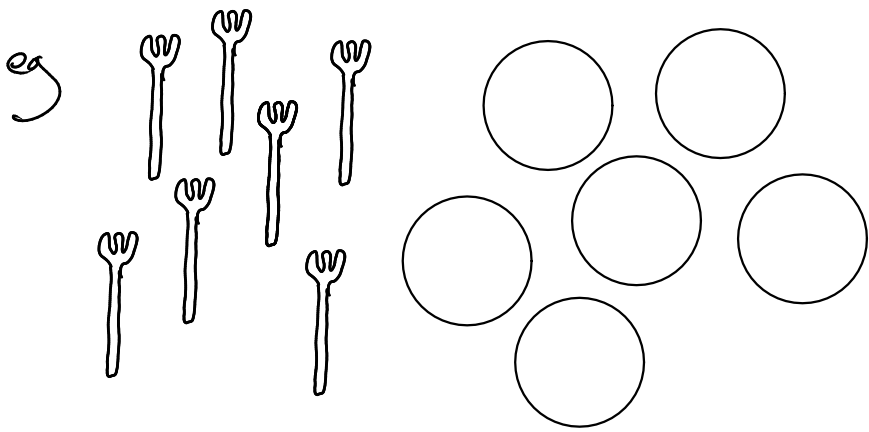
yes $4 = \{0, 1, 2, 3\}$

let $f = \{(0, 2), (1, 3), (2, 8), (3, 12)\}$

then f is one-to-one and onto

this is what a **one-to-one correspondence** is

The **idea** here is that if I want to know if two sets have the **same** number of elements then I don't have to know how to count the elements, I just have to **match** them up and see if I get any **leftovers**.



For the break

find an example of a set which is equivalent
to a **proper subset** of itself

eg $f: \omega \rightarrow \omega$
 $f(n) = n^+$ or $f: \omega \rightarrow \omega$
 $f(n) = n^2$ etc.

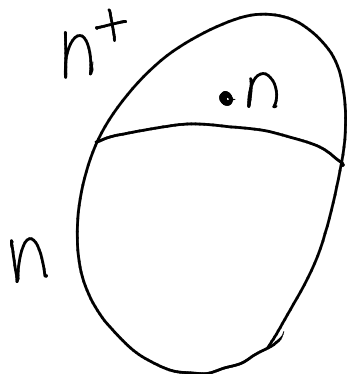
Fortunately this doesn't happen for individual natural numbers

Claim If $n \in \omega$ then n is not equivalent to a proper subset of n

proof by induction: let S be the set of $n \in \omega$ which are not equivalent to any proper subset of themselves.

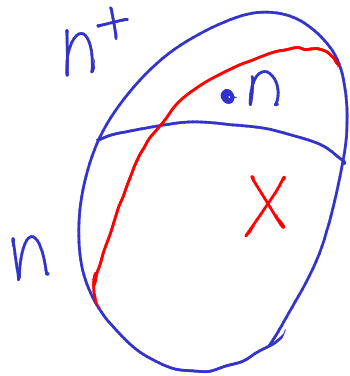
$\emptyset \in S$ since \emptyset has no proper subset

Take $n \in S$ and say $f: n^+ \rightarrow x$ is one-to-one and onto with $x \subsetneq n^+$



But $n^+ = n \cup \{n\}$

if $n \in X$

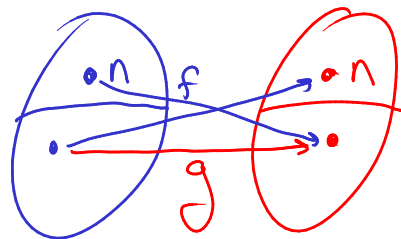


then $X - \{n\} \subsetneq n$

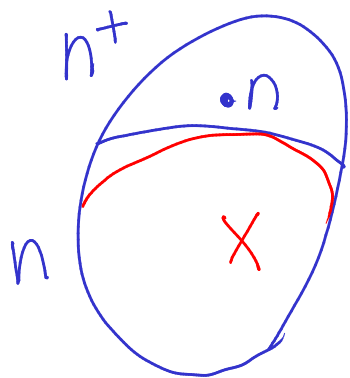
so we just need to find
a one-to-one, onto map
 $g: n \rightarrow X - \{n\}$ to get
a contradiction to $n \in S$

if $f(n) = n$ let g be f restricted
to $n^+ - \{n\} = n$

if $f(n) \neq n$ let $g(a) = f(a)$ for $f(a) \neq n$
and $g(a) = f(n)$ for $f(a) = n$



if $n \notin X$



then let g be f
restricted to $n^+ - \{n\} = n$

then g maps onto $X - \{f(n)\}$

as $X - \{f(n)\} \subsetneq n$

again contradicting $n \in S$

all together we get that if $n \in S$ then $n+1 \in S$
so by the principle of mathematical induction $S = \omega$
and so no natural number is equivalent to a proper
subset of itself.

Claim A set can be equivalent to at most one natural
number

proof Suppose S is equivalent to n and to m
with $n \neq m$, $n, m \in \omega$.

Then n is equivalent to m (can you explain why?)
But either $n \subsetneq m$ or $m \subsetneq n$
and so this contradicts the previous claim.

Now we can **define** a set A to be **finite** if it is equivalent to some natural number, and **infinite** otherwise

Also **define** the **size** or **number of elements** of a finite set A to be the unique natural number equivalent to A

use the notation **$\#A$** for the size of A

This notion of size corresponds to our usual notion of size

for example

if $A \subseteq B$ then $\#A \leq \#B$

proof Say B is equivalent to n
by the map f

Then f restricted to A gives
 A equivalent to a subset X
of n . Suppose $\#X > n$

Then we would have a

one-to-one correspondence between $\#X$ and X , but $X \subseteq \mathbb{N} \subsetneq \#X$ so we would have a one-to-one correspondence between $\#X$ and a proper subset of $\#X$, contradiction

③ Next time

- Summary of our axioms so far and outlook
- The axiom of choice

Please read Halmos Section 15