

Math 821, Spring 2013, Lecture 12

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1 Duals

Definition. (1) Let V be a finite dimensional vector space over k , then $V^* = \text{Hom}(V, k)$ the space of linear maps from V to k .

(2) If $\Phi : V \rightarrow W$ is a linear map, then $\Phi^* : W^* \rightarrow V^*$ by $(\Phi^*(f))(v) = f(\Phi(v))$.

(3) If v_1, v_2, \dots, v_n is a basis for V , let $f_i \in V^*$ be the map $f_i(v_i) = 1$, $f_i(v_j) = 0$ for $i \neq j$. Then f_1, f_2, \dots, f_n is a basis of V^* , called the dual basis.

Note. (2) says taking duals reverses arrows. So we should expect algebra \rightarrow coalgebra and vice versa.

Definition. A graded k -vector space $V = \bigoplus_{i=0}^{\infty} V_i$ is of finite type if each V_i is finite dimensional. Note if \mathcal{C} is a combinatorial class then $V\mathcal{C} = \bigoplus_{n=0}^{\infty} V\mathcal{C}_n$ is of finite type.

Definition. Let $V = \bigoplus_{n=0}^{\infty} V_n$ be a graded vector space of finite type. Then the restricted dual is $V^o = \bigoplus_{n=0}^{\infty} V_n^*$.

Note. The elements of V^o are linear maps from $V \rightarrow k$ which vanish on all but finitely many of the V_n .

So if we have a graded algebra, take its restricted dual (from now on just call this the dual) get a coalgebra and vice versa. So a graded bialgebra will have a dual which is also a graded bialgebra.

Note. *Connected is preserved under duals because $k^* = \text{Hom}(k, k) \cong k$. So the dual of a graded connected finite type Hopf algebra is a graded connected finite type Hopf algebra.*

But what does this look like concretely?

Let A be a graded connected finite type Hopf algebra. What is Δ_{A° ?

$$\begin{aligned}\Delta_{A^\circ} : A^\circ &\rightarrow A^\circ \otimes A^\circ \\ \Delta_{A^\circ}(f)(a \otimes b) &= f(ab)\end{aligned}$$

where $A^\circ = \bigoplus_{n=0}^{\infty} \text{Hom}(A_n, k) \subseteq \text{Hom}(A, k)$. What is \cdot_{A° ?

$$\begin{aligned}\cdot_{A^\circ} : A^\circ \otimes A^\circ &\rightarrow A^\circ \\ (\cdot_{A^\circ}(f \otimes g))(a) &= (f \otimes g)(\Delta(a))\end{aligned}$$

Write this in terms of a basis.

Proposition 1. *Say $\{a_i\}_{i \in I}$ a basis for A a graded connected finite type Hopf algebra and let $\{f_i\}_{i \in I}$ be the dual basis, write*

$$\begin{aligned}a_j a_k &= \sum_{i \in I} c_{j,k}^i a_i \\ \text{then } \Delta_{A^\circ}(f_i) &= \sum_{(j,k) \in I \times I} c_{j,k}^i f_j \otimes f_k\end{aligned}$$

and dually, write

$$\begin{aligned}\Delta(a_i) &= \sum_{(j,k) \in I \times I} d_{j,k}^i a_j \otimes a_k \\ \text{then } f_j \cdot_{A^\circ} f_k &= \sum_{i \in I} d_{j,k}^i f_i\end{aligned}$$

Proof. It suffices to prove the first part by the previous observation.

$$\Delta_{A^\circ}(f_i)(a_j \otimes a_k) = f_i(a_j a_k) = f_i\left(\sum_{l \in I} c_{j,k}^l a_l\right) = \sum_{l \in I} c_{j,k}^l f_i(a_l) = c_{j,k}^i$$

□

Example. Let $TV = \text{words}$ let $A = TV$, what is A° ?

A basis for A is given by words, so the dual basis for A° is indexed by words, thinking of a basis element as its index, we can view A° as also being a Hopf algebra of words.

Recall multiplication on A is concatenation and comultiplication on A is anti-shuffle. So comultiplication of A° is deconcatenation

$$\Delta_{A^\circ}(abcd) = \mathbb{1} \otimes abcd + a \otimes bcd + ab \otimes cd + abc \otimes d + abcd \otimes \mathbb{1}$$

multiplication on A° is shuffle

$$ab \cdot_{A^\circ} cd = abcd + acbd + cabd + acdb + cadb + cdab$$

Example. Let H be the Connes-Kreimer Hopf algebra of rooted trees, what is H° ?

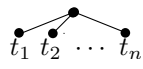
What is multiplication in H° , it has to be dual to taking admissible cuts so it will be a grafting operator. eg.

$$\text{tree}_1 \cdot_{H^\circ} \text{tree}_2 = \text{tree}_1 \text{ grafted to tree}_2 + \text{tree}_2 \text{ grafted to tree}_1 + \text{tree}_1 \text{ and tree}_2 \text{ grafted to a new root}$$

The comultiplication in H° is the dual of disjoint union

$$\begin{aligned} \Delta_{H^\circ}(\text{tree}_1 \text{ and } \text{tree}_2) = & \text{tree}_1 \otimes \text{tree}_2 + \mathbb{1} \otimes \text{tree}_1 + \text{tree}_2 \otimes \mathbb{1} + \text{tree}_1 \otimes \text{tree}_2 + \text{tree}_2 \otimes \text{tree}_1 \\ & + \text{tree}_1 \otimes \text{tree}_2 + \text{tree}_2 \otimes \text{tree}_1 + \text{tree}_1 \otimes \text{tree}_2 + \text{tree}_2 \otimes \text{tree}_1 \\ & + \text{tree}_1 \otimes \text{tree}_2 + \text{tree}_2 \otimes \text{tree}_1 \end{aligned}$$

I partition the connected components into two parts in all possible ways.

This is almost isomorphic to the Grossman-Larson Hopf algebra in the following way for each forest t_1, t_2, \dots, t_n of H° form the tree , these are the basis for the Grossman-Larson Hopf algebra G (note no empty tree here). Now the size of a tree is its number of edges, so \bullet is size 0.

If we define the dual space of \mathcal{H} with respect to a symmetric factor, i.e,

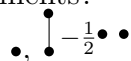
$$\begin{aligned} \Delta_{H^\circ}(\text{tree}_1 \text{ and } \text{tree}_2) = & \text{tree}_1 \otimes \text{tree}_2 + \mathbb{1} \otimes \text{tree}_1 + \text{tree}_2 \otimes \mathbb{1} + 2 \text{tree}_1 \otimes \text{tree}_2 \\ & + 2 \text{tree}_2 \otimes \text{tree}_1 + \text{tree}_1 \otimes \text{tree}_2 + \text{tree}_2 \otimes \text{tree}_1 + 2 \text{tree}_1 \otimes \text{tree}_2 \\ & + 2 \text{tree}_2 \otimes \text{tree}_1 + \text{tree}_1 \otimes \text{tree}_2 \end{aligned}$$

We can import \cdot_{H° , Δ_{H° . $t_1 \cdot_G t_2$ grafts the child of the root of t_1 into t_2 (actually usually define the opposite of this, i.e. graft children of t_2 into t_1) and $\Delta_G(t)$ partitions the children of root of t onto the two sides of \otimes and given them a new root on each side.

References. Victor Reiner, *Hopf Algebras In Combinatorics, Cha 1.*
 For Grossman-Larson Hopf algebra: see arXiv:math/0003074, Florin Panaite, *Relating the Connes-Kreimer and Grossman-Larson Hopf algebras built on rooted trees.*
 arXiv:math/0201253, Michael E. Hoffman, *Combinatorics of Rooted Trees and Hopf Algebras.*

2 B_+

Working in the Connes-Kreimer Hopf algebra H , what are the primitive elements?



(recall primitive $\Delta(a) = a \otimes \mathbb{1} + \mathbb{1} \otimes a$)

$$\text{Check } \tilde{\Delta}(\text{root node}) = \bullet \otimes \bullet - \frac{1}{2}(\bullet \otimes \bullet + \bullet \otimes \bullet) = 0$$

Definition. $B_+ : H \rightarrow H$ is the linear function which takes a forest t_1, t_2, \dots, t_n



and returns $t_1 t_2 \dots t_n$

What is ΔB_+ ?

$$\begin{aligned} \Delta B_+(t_1 t_1 \dots t_n) &= \Delta(\text{root node}) \\ &= t_1 t_2 \dots t_n \otimes \mathbb{1} + (id \otimes B_+) \prod_{i=1}^n \sum_{\substack{c \\ \text{admissible cut of } t_i}} P_c(t_i) \otimes R_c(t_i) \\ &= B_+(t_1 t_1 \dots t_n) \otimes \mathbb{1} + (id \otimes B_+)(\Delta(t_1 t_1 \dots t_n)) \\ \Rightarrow \Delta B_+ &= B_+ \otimes \mathbb{1} + (id \otimes B_+) \Delta \end{aligned}$$

3 3 line summary of cohomology

- You need a family maps b_n from objects of size n to objects of size $n+1$ with $b^2 = b_{n+1}b_n = 0$.
- Take quotients $Ker(b)/Im(b)$.
- Use these to understand your original objects.

For us we want "objects of size n" to be $Hom(H, H^{\otimes n})$ (actually H could be any bialgebra here) and $b : Hom(H, H^{\otimes n}) \rightarrow Hom(H, H^{\otimes n+1})$.

$$bL = (id \otimes L)\Delta + \sum_{i=1}^n (-1)^i \Delta_i L + (-1)^{n+1} L \otimes \mathbb{1}$$

where, $\Delta_i = id \otimes \dots \otimes \Delta \otimes id \dots \otimes id$, Δ is the i -th part.

This gives the Hochschild cohomology of bialgebras. If I were to do this, one of the first thing I'd need to know is $Ker(b_1)$

$$\begin{aligned} 0 &= b_1 L \\ &= (id \otimes L)\Delta - \Delta L + L \otimes \mathbb{1} \\ \text{so } \Delta L &= L \otimes \mathbb{1} + (id \otimes L)\Delta \end{aligned}$$

that's the property B_+ has, this is called the 1-cocycle property.

Comment. This 1-cocycle property is really important in these Renormalization Hopf algebras (like Connes-Kreimer) but I haven't seen it appear in other combinatorial hopf algebras.

4 Specifications and Combinatorial Dyson-Schwinger equations

Again let H be the Connes-Kreimer Hopf algebra. I can use B_+ to give combinatorial specifications in a different languages.

Example. $T(x) = \mathbb{1} - xB_+(\frac{1}{T(x)})$

I want a solution to this in $H[[x]]$, expand this recursively

$$T(x) = \mathbb{1} - x \bullet - x^2 \begin{array}{c} \bullet \\ | \\ \bullet \end{array} - x^3 \left(\begin{array}{c} \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \diagdown \\ \bullet \end{array} \right) - x^4 \left(\begin{array}{c} \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \diagdown \\ \bullet \end{array} + 2 \begin{array}{c} \bullet \\ \diagup \diagdown \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \diagdown \\ \bullet \end{array} \right) + O(x^5)$$

this is just $\mathcal{T} = \mathcal{E} + \mathcal{Z} \times Seq(\mathcal{T} - \mathcal{E})$, plane rooted trees then forget the plane structure giving above coefficients.

Example. $T(x) = \mathbb{1} + xB_+(T(x)^2)$

$$T(x) = \mathbb{1} + x \bullet + x^2 \begin{array}{c} \bullet \\ | \\ \bullet \end{array} + x^3 \left(4 \begin{array}{c} \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \diagdown \\ \bullet \end{array} \right) + O(x^4)$$

$\mathcal{T} = \mathcal{E} + \mathcal{Z} \times \mathcal{T}^2$ and forget the extra structure to get the coefficients.

These equations (and more general ones) are called combinatorial Dyson-Schwinger equations (this name is given by Dr. Karen Yeats) they give some specifications in the Hopf algebra context. They are physically important and give the sums of trees/ feynman graphs which contributes to a given physical process.